Generating Posets beyond **N**

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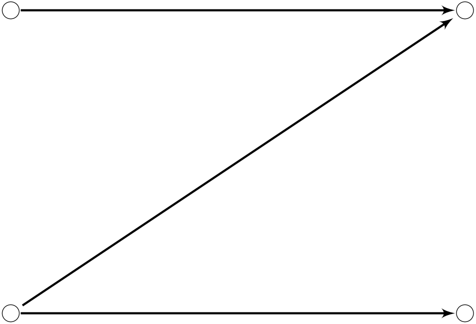
RAMiCS 2020

Motivation



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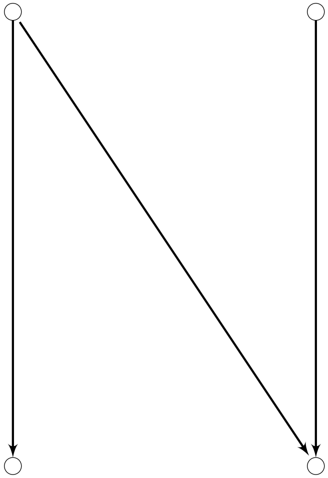

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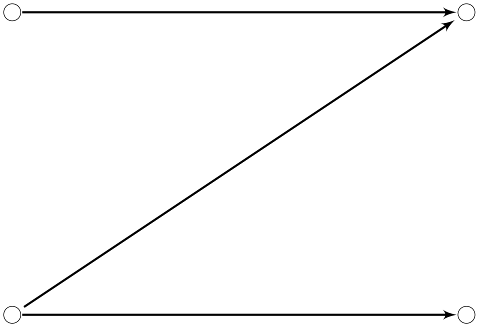

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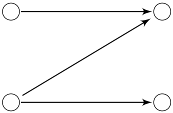
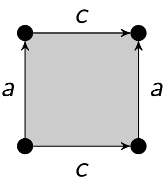
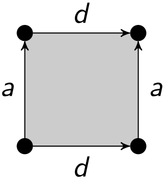


- Kleene algebra is nice and useful
  - also its extensions: semimodules, tests, domain, ...
- Concurrent Kleene algebra: extension of KA for concurrency
  - [Hoare, Möller, O'Hearn, Struth, van Staden, Villard, Wehrman, Zhu '09, '11, '16]
- Kleene algebra plus parallel composition
- the free CKA (minus some details): sets of series-parallel pomsets
  - labeled posets with concatenation & parallel composition
- *Something's amiss in concurrent Kleene algebra*

Example, Using Higher-Dimensional Automata





$\begin{pmatrix} a \\ c \end{pmatrix}$ $\begin{pmatrix} a \\ d \end{pmatrix}$

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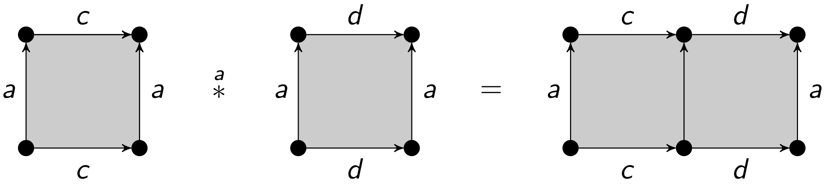

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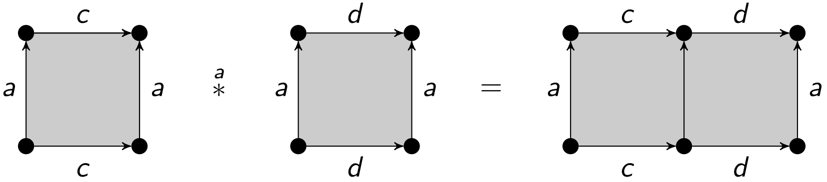

$$\begin{pmatrix} a \\ c \end{pmatrix} \qquad \begin{pmatrix} a \\ d \end{pmatrix} \qquad \begin{pmatrix} & a & \\ c & \longrightarrow & d \end{pmatrix}$$

$$\begin{pmatrix} a \\ c \end{pmatrix} \Big\| \begin{pmatrix} a \\ d \end{pmatrix} = \begin{pmatrix} a \\ c \\ a \\ d \end{pmatrix} \text{ ??}$$

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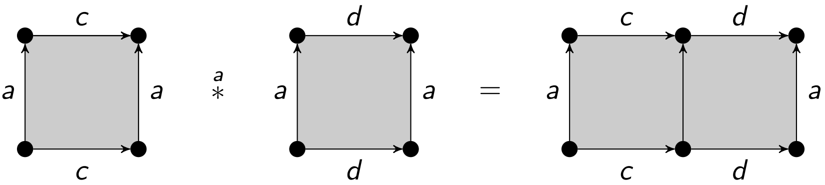
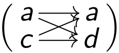


$\begin{pmatrix} a \\ c \end{pmatrix}$ $\begin{pmatrix} a \\ d \end{pmatrix}$ $\left(c \xrightarrow{a} d \right)$

$\begin{pmatrix} a \\ c \end{pmatrix} * \begin{pmatrix} a \\ d \end{pmatrix} =$  ??

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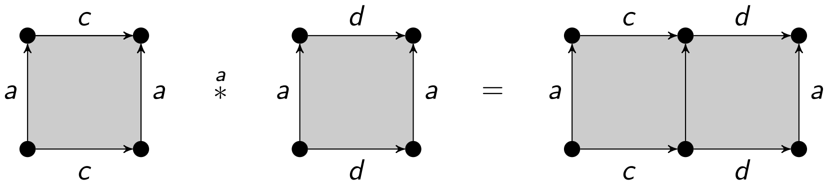
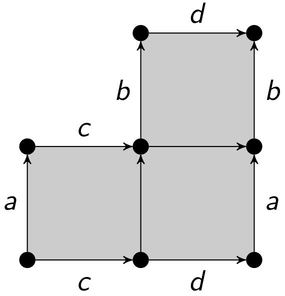


$$\begin{pmatrix} a \\ c \end{pmatrix} \overset{a}{*} \begin{pmatrix} a \\ d \end{pmatrix} = \begin{pmatrix} & a & \\ c & \longrightarrow & d \end{pmatrix}$$

- new gluing operation on pomsets, to *continue events across compositions*

Another Example



$$\begin{pmatrix} a \\ c \end{pmatrix} \overset{a}{*} \begin{pmatrix} a \\ d \end{pmatrix} \overset{d}{*} \begin{pmatrix} b \\ d \end{pmatrix} = \begin{pmatrix} a \longrightarrow b \\ \nearrow \\ c \longrightarrow d \end{pmatrix}$$

- this is the **N** pomset, which is not series-parallel
- hence our title, Generating Posets beyond **N**

Series-Parallel Posets

- a poset: *finite* set P plus partial order $\leq$: reflexive, transitive, antisymmetric
- parallel composition of posets $(P_1, \leq_1)$, $(P_2, \leq_2)$:

$$P_1 \otimes P_2 = (P_1 \sqcup P_2, \leq_1 \cup \leq_2)$$

$\upuparrows$ disjoint union

- serial composition:

$$P_1 * P_2 = (P_1 \sqcup P_2, \leq_1 \cup \leq_2 \cup P_1 \times P_2)$$

$\upuparrows$ P_1 before P_2

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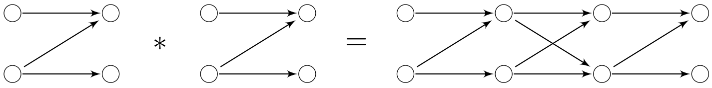


Series-Parallel Posets

Definition (Winkowski '77, Grabowski '81)

A poset is series-parallel (sp) if it is empty or can be obtained from the singleton poset by a finite number of serial and parallel compositions.

Theorem (Grabowski '81)

A poset is sp iff it does not contain ***N*** *as an induced subposet.*

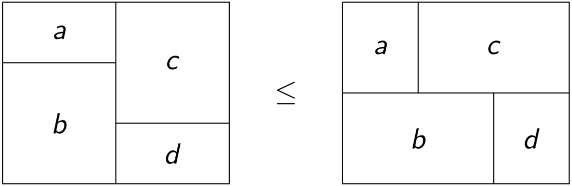
The equational theory of sp-posets is well-understood: [Gischer '88], [Bloom-Esik '96]

Concurrent Monoids

Definition (Gischer '88, Hoare *et al.* '11)

A concurrent monoid is an ordered bimonoid $(S, \leq, *, \|, 1)$ with shared $*$-$\|$-unit 1 which satisfies weak interchange:

$$(a\|b) * (c\|d) \leq (a * c)\|(b * d)$$



- subsumption on posets: $P \preceq Q$ if P "has more order" than Q

Theorem (Gischer '88, Bloom-Esik '96)

The set of sp-posets under subsumption is the free concurrent monoid.

Problem & Solution

- we like the **N** poset, but it's not series-parallel
- in fact, **N**'s are everywhere: for example, *producer-consumer:*

$$
\begin{array}{ccccccccc}
p_1 & \longrightarrow & p_2 & \longrightarrow & p_3 & \longrightarrow & p_4 & \longrightarrow & \cdots \\
\downarrow & & \downarrow & & \downarrow & & \downarrow & & \\
c_1 & \longrightarrow & c_2 & \longrightarrow & c_3 & \longrightarrow & c_4 & \longrightarrow & \cdots
\end{array}
$$

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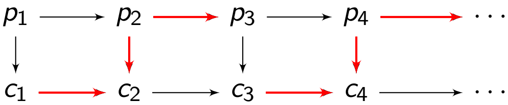

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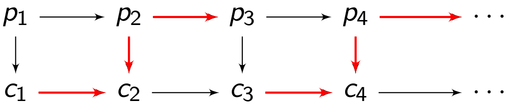


Problem

Find a class of posets which includes ***N*** *(and sp-posets) and which has good algebraic properties.*

Our Proposal

Posets with interfaces with parallel and gluing composition.


Neuer Weg
Alter Weg

Posets with Interfaces

Definition

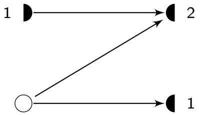
A poset with interfaces (iposet) is a poset P plus two injections

$$[n] \xrightarrow{s} P \xleftarrow{t} [m]$$

such that $s[n]$ is minimal and $t[m]$ is maximal in P.

- ($[n] = \{1, \dots, n\}$; $S \subseteq P$ *minimal* if $p \not< s$ for all $p \in P$, $s \in S$)
- (there are 25 non-isomorphic iposets with underlying $\mathbf{N}$)

Interfaces



Def.: Iposet $s : [n] \to P \leftarrow [m] : t$; $s[n] \subseteq P_{\min}$, $t[m] \subseteq P_{\max}$.

- s: starting interface ; t: terminating interface
- events in $t[m]$ are *unfinished* ; events in $s[n]$ are "*unstarted*"

Definition

The gluing composition of iposets $s_1 : [n] \to (P_1, \leq_1) \leftarrow [m] : t_1$ and $s_2 : [m] \to (P_2, \leq_2) \leftarrow [k] : t_2$:

$$P_1 * P_2 = \begin{cases} (P_1 \sqcup P_2)/t_1(i) = s_2(i) \\ \leq_1 \cup \leq_2 \cup (P_1 \setminus t_1[m]) \times (P_2 \setminus s_2[m]) \end{cases}$$

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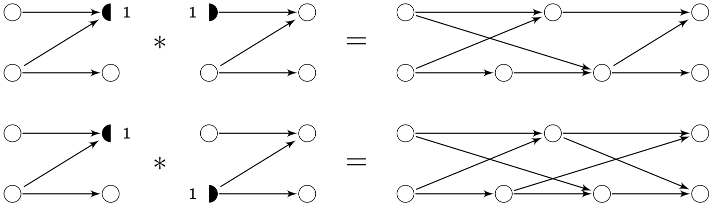

- only defined if terminating int. of P_1 is equal to starting int. of P_2
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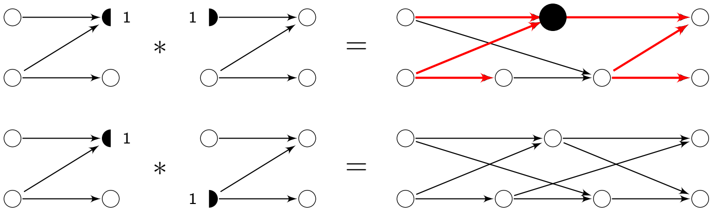

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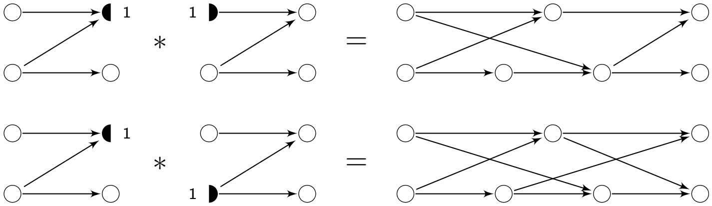

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Parallel Composition

- parallel composition of iposets: put posets in parallel and renumber interfaces
- for $[n_1] \to P_1 \leftarrow [m_1]$ and $[n_2] \to P_2 \leftarrow [m_2]$, have $[n_1 + n_2] \to P_1 \otimes P_2 \leftarrow [m_1 + m_2]$
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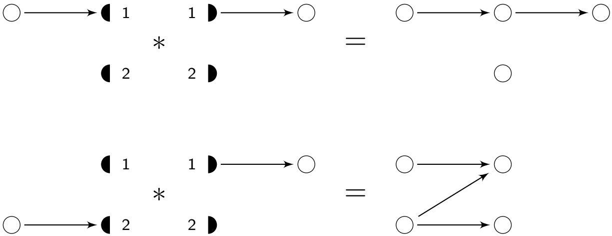
$\upuparrows$



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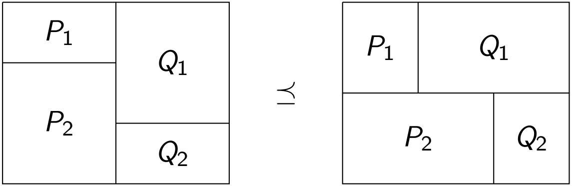
$\Uparrow$

$$(P_1 \otimes P_2) * (Q_1 \otimes Q_2) \preceq (P_1 * Q_1) \otimes (P_2 * Q_2)$$



Gluing-Parallel Iposets

- the four singletons:

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  ○   1 ◗   ◖ 1   1 ⧓ 1

- recall *sp-posets*: generated from ○ using $*$ and $\otimes$
  - sp-posets are *freely* generated
  - P is sp iff P is **N**-free

Gluing-Parallel Iposets

- the four singletons:

  ◯ 1◗ ◖1 1⧓1

- recall *sp-posets*: generated from ◯ using $*$ and $\otimes$
  - sp-posets are *freely* generated
  - P is sp iff P is **N**-free
- gp-iposets: generated from ◯ , 1◗ , ◖1 , 1⧓1 using $*$ and $\otimes$

Proposition

Gp-iposets are freely generated, except for the relations

$$\begin{pmatrix} \text{◖}1 \\ P \end{pmatrix} * \begin{pmatrix} 1\text{◗} \\ Q \end{pmatrix} = \begin{pmatrix} \text{◯} \\ P * Q \end{pmatrix} \qquad \begin{pmatrix} \text{◖}1 \\ P \end{pmatrix} * \begin{pmatrix} 1\text{⧓}1 \\ Q \end{pmatrix} = \begin{pmatrix} \text{◖}1 \\ P * Q \end{pmatrix}$$

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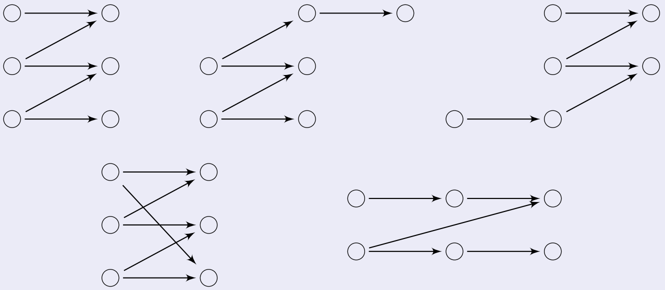
Forbidden Substructures

Proposition

If P is gp, then it does not contain any of the following as induced subposets:



- unlike for *sp*-posets, that's not an iff (we don't know)
- but these five are the only posets on ≤ 6 points which are not gp

Some Counting, up to Isomorphism

n	P(n)	SP(n)	GP(n)	IP(n)	GPI(n)
0	1	1	1	1	1
1	1	1	1	4	4
2	2	2	2	17	16
3	5	5	5	86	74
4	16	15	16	532	419
5	63	48	63	???	2980
6	318	167	313	???	26566
7	2045	602	???	???	???
OEIS	A000112	A003430	A079566 ?	n.a.	n.a.

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- slow Python implementation
- bottleneck is isomorphism checking

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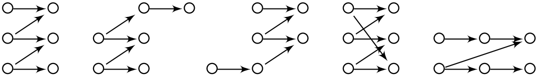
- $\mathsf{P}(6) - \mathsf{GP}(6) = 5$: 

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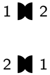
- the only iposet on 2 points which is not gp:



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- Online Encyclopedia of Integer Sequences A079566: the number of connected graphs without induced sub-C_4

Conclusion

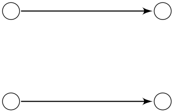
- posets with interfaces for concurrency
- instead of concurrent monoid, small category with lax tensor
  - a "multi-object concurrent monoid"
- gluing-parallel iposets include sp-posets and the **N**
  - they also include all interval orders
  - **I I**-free; useful in concurrency [Wiener 1914], [Lamport '86], [van Glabbeek '90], [Vogler '91], [Janicky '93], etc.
- generation is "almost free"
- characterization by forbidden substructures?

Ongoing and Future Work

- Concurrent Kleene algebra:
  - concurrent monoid $\rightsquigarrow$ concurrent semiring
  - multi-object concurrent monoid $\rightsquigarrow$ bicategories with lax tensors?
  - and the stars?
  - relation to synchronous Kleene algebra?
- CKA with domain:
  - domain elements are "structure-less" iposets
  - relation to higher-dimensional modal logic?
  - higher-dimensional modal Kleene algebra
- Languages of higher-dimensional automata:
  - sets of interval orders
  - concatenation of HDA $\approx$ gluing of (sets of) interval orders
  - theory of regular languages for concurrency?

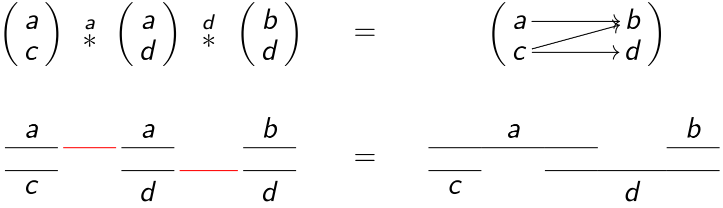
6 Posets for Concurrency: Interval Orders

Posets for Concurrency: Interval Orders

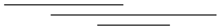


- posets which are good for concurrency?
- already in [Wiener 1914], then [Winkowski '77], [Lamport '86], [van Glabbeek '90], [Vogler '91], [Janicky '93], etc.
- interval orders: posets which have representation as (real) intervals, ordered by $\max_1 \le \min_2$
- Lemma (Fishburn '70): A poset is interval iff it does not contain $\mathbf{II} = \left(\begin{smallmatrix} \cdot \longrightarrow \cdot \\ \cdot \longrightarrow \cdot \end{smallmatrix} \right)$ as induced subposet.
- intuitively: if $a \longrightarrow b$ and $c \longrightarrow d$, then also $a \longrightarrow d$ or $c \longrightarrow b$

Gluing of Interval Orders



Interval Orders vs ST-Traces

- An ST-trace: $a^+ \, b^+ \, a^+ \, a^- \, a^- \, b^-$ [van Glabbeek '90]
- as intervals: 

Proposition

ST-traces up to the equivalence generated by $a^+ b^+ \sim b^+ a^+$ and $a^- b^- \sim b^- a^-$ are the same as labeled interval orders.